Flow networks
COMS20010 (Algorithms II)

John Lapinskas, University of Bristol

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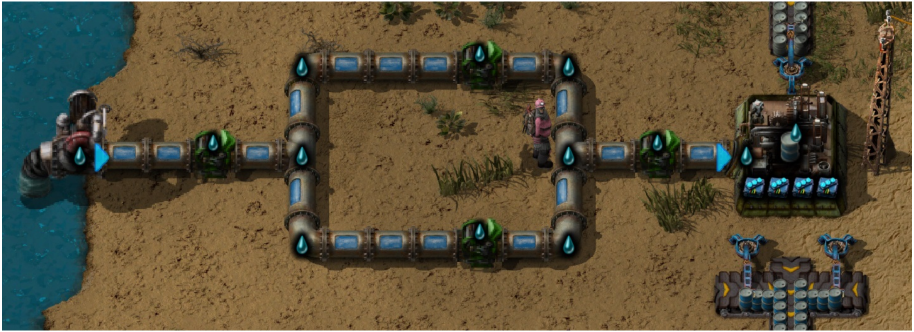
- Survey design;
- Professional baseball. (See KT 7.12!)

For now, let's just consider a toy problem. One pump supplies water for one factory, passing through a network of pipes of different capacities.



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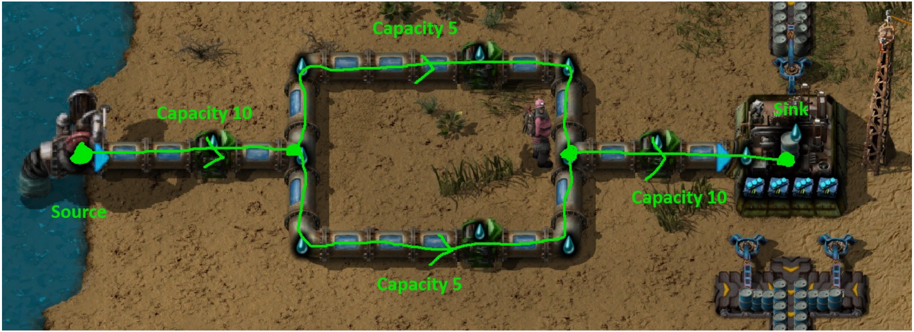

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The **value** of f, denoted $v(f)$, is $f^+(s)$.

The problem: Find a **maximum flow**: a flow f maximising $v(f)$.

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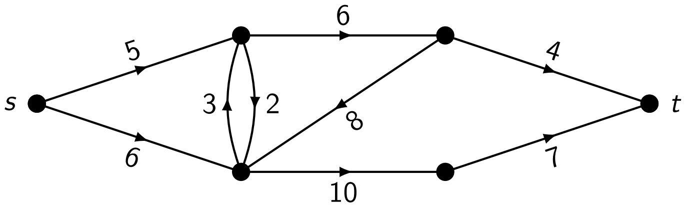

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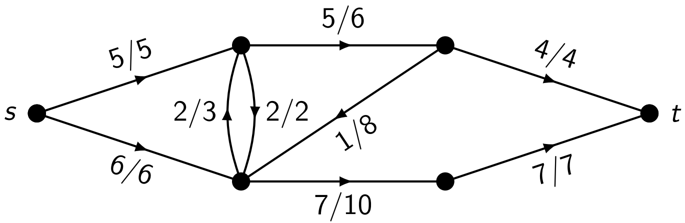

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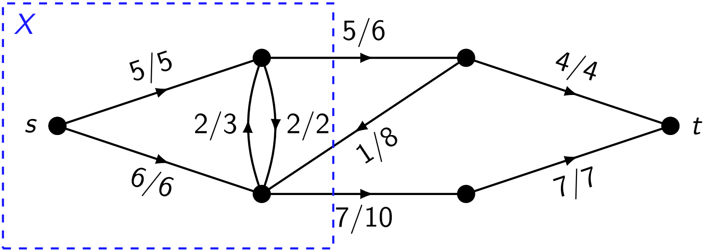

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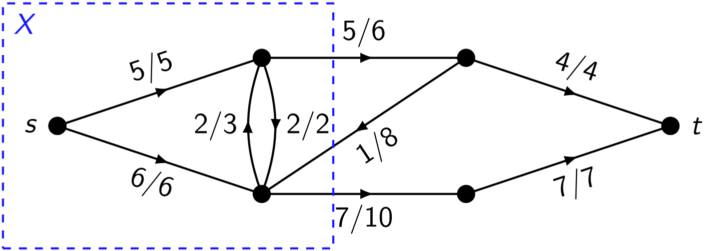

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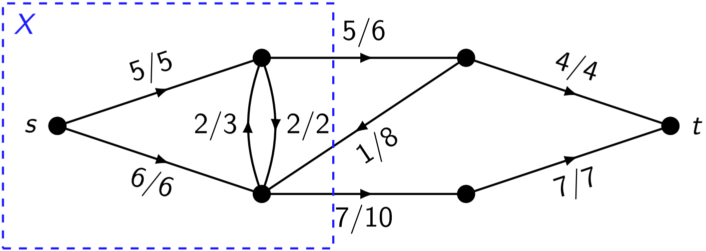

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For example, here $f^+(X) = 5 + 7 = 12$ and $f^-(X) = 1$.

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For all $e \subseteq X$, $f(e)$ appears once on each side; after cancelling those terms we're left with $f^+(X) = f^-(X)$. □

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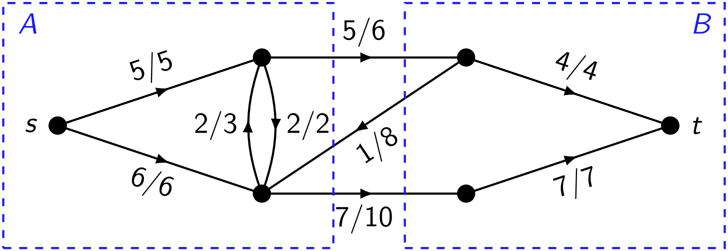

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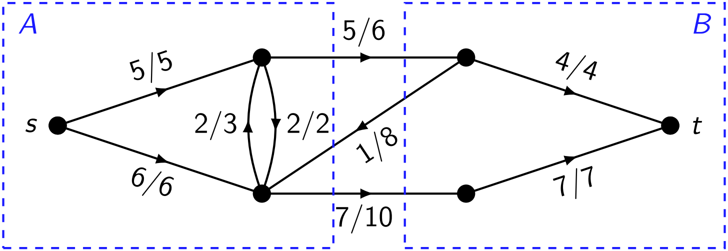

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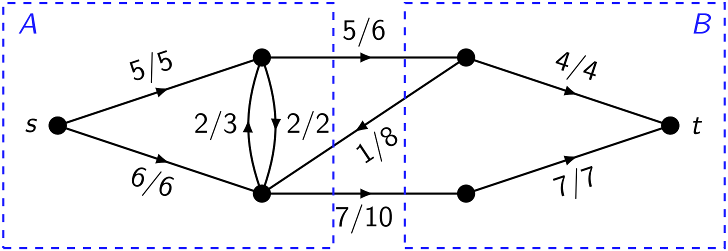

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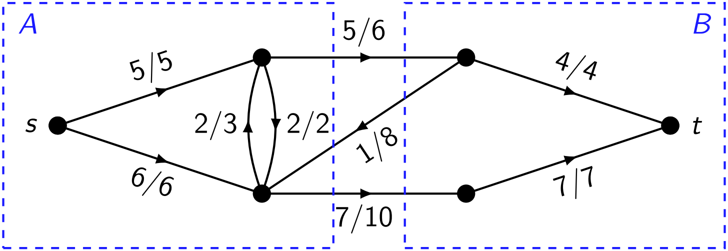

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But $f^+(A \setminus \{s\}) = f^+(A) - f(s, B)$ and $f^-(A \setminus \{s\}) = f^-(A) + f(s, A)$...

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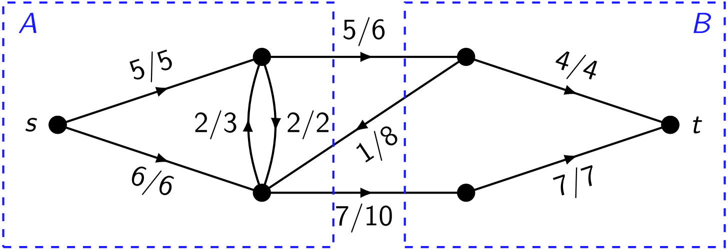

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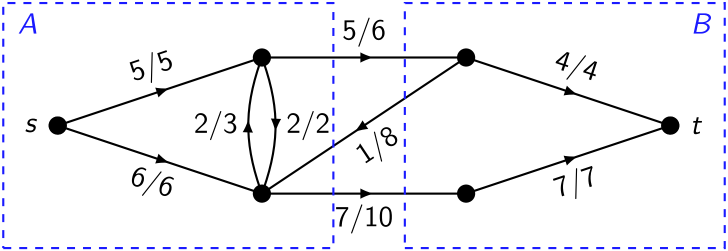

Lemma 2: For all cuts (A, B), $f^+(A) - f^-(A) = f^-(B) - f^+(B) = v(f)$.

Proof: By Lemma 1, we have $f^+(A \setminus \{s\}) = f^-(A \setminus \{s\})$.

Rearranging $f^+(A) - f(s, B) = f^-(A) + f(s, A)$:

The **value** of a flow f, denoted $v(f)$, is $f^+(s)$.
We write $f^+(A) := \sum_{e \text{ out of } A} f(e)$ and $f^-(A) := \sum_{e \text{ into } A} f(e)$.

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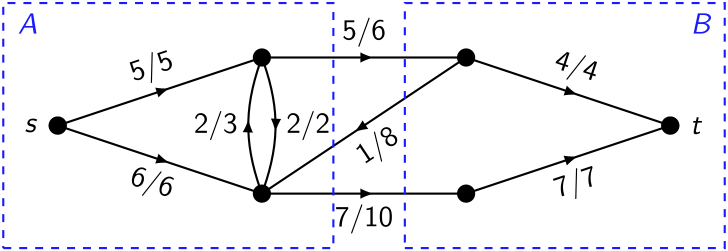

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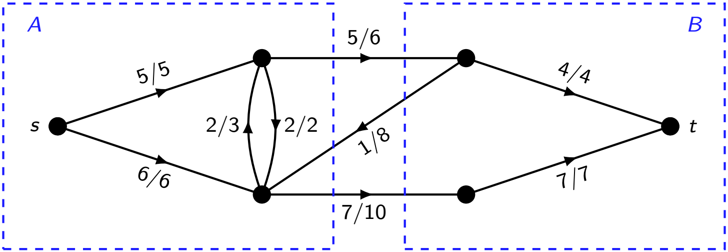

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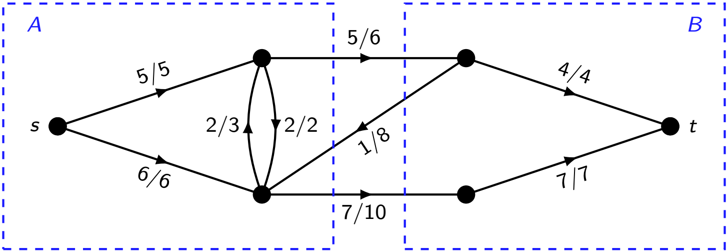

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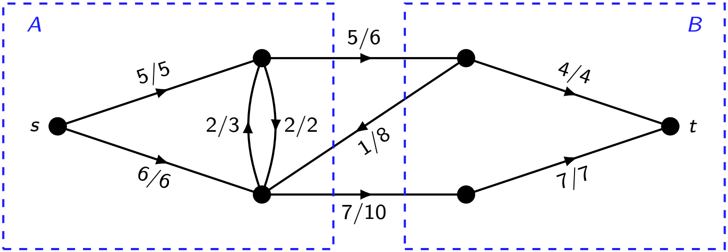

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Lemma 2 implies we could have defined $v(f)$ via **any** cut in the network. In particular, $f^+(s) = f^-(t)$.